

Lecture 1 - Basics of Quantum Information

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What is Quantum Information Science?

Quantum mechanics + Information theory + computer Science

≡ info stored in and manipulation of quantum states

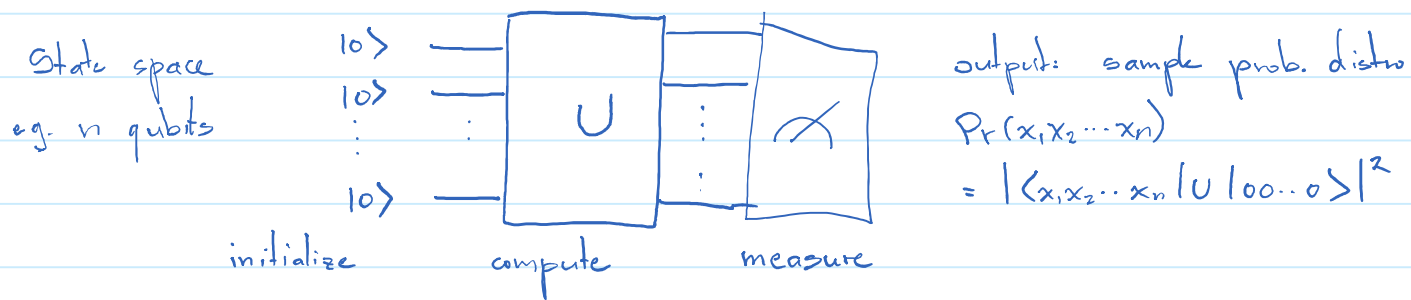
Classical info: bit strings $x_1 x_2 \dots x_n$ $x_i \in \{0, 1\}$

Quantum info: qubit strings $\sum_{x_1=0,1} \dots \sum_{x_n=0,1} c_{x_1 \dots x_n} |x_1\rangle \otimes \dots \otimes |x_n\rangle$

Q: How is quantum info different?

- true randomness
- uncertainty
- no cloning
- entanglement
- superpositions ("quantum parallelism")

Q: What is a quantum computation?



→ hopefully, answer x to an interesting problem occurs w/ high prob.!

- won't discuss computation & algorithms in these lectures
- instead, mathematical properties of QI & how it evolves, w/ some applications to HEP
- central idea: Quantum Channel
→ most general map b/w quantum states

- Outline:
- ① Basics of QI (states, entropy, black holes)
 - ② Exercises: relative entropy & distinguishability
 - ③ Channels
 - ④ Quantum Error Correction
 - ⑤ Holography as QEC



States

Let \mathcal{H} : Hilbert space, $\dim \mathcal{H} = d$, ON basis $\{|i\rangle\}_{i=1}^d$ ↖ countable ∞ or

- generic pure state $|\psi\rangle \in \mathcal{H}$

$$|\psi\rangle = \sum_{i=1}^d c_i |i\rangle \quad \langle \psi | \psi \rangle = \sum |c_i|^2 = 1$$

- generic mixed state $\rho \in \mathcal{L}(\mathcal{H})$ density matrix

$$\rho = \sum_{i,j} \rho_{ij} |i\rangle\langle j|$$

Properties of density matrices:

1. $\rho = \rho^\dagger$ (Hermitian)
2. $\langle \psi | \rho | \psi \rangle \geq 0$ (pos. semidef)
3. $\text{Tr} \rho = 1$ (normalization)
4. \exists basis $\{|a\rangle\}$: $\rho = \sum_{a=1}^d p_a |a\rangle\langle a|$, $p_a \geq 0$, $\sum p_a = 1$
5. If only one $p_a \neq 0 \rightarrow \rho = |a\rangle\langle a|$ pure
6. $\langle 0 | \rho | 0 \rangle = \text{Tr}(\rho |0\rangle\langle 0|)$

Notation: $\mathcal{S}(\mathcal{H}) \equiv$ set of density matrices on \mathcal{H}

Composition of Hilbert spaces

• $\mathcal{H}_A = \text{span}\{|i\rangle_A\}_{i=1}^{d_A}$ $\mathcal{H}_B = \text{span}\{|\mu\rangle_B\}_{\mu=1}^{d_B}$

$$\mathcal{H}_A \otimes \mathcal{H}_B \equiv \text{span}\{|i\rangle_A \otimes |\mu\rangle_B\}_{i=1, \mu=1}^{d_A, d_B}$$

$$\dim(\mathcal{H}_A \otimes \mathcal{H}_B) = d_A d_B$$

$$|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \Rightarrow |\psi\rangle = \sum_{i=1}^{d_A} \sum_{\mu=1}^{d_B} c_{i\mu} |i\rangle_A \otimes |\mu\rangle_B$$

Defⁿ $\langle \mu | \rho : \mathcal{H}_{AB} \rightarrow \mathcal{H}_A \quad \langle \mu | \rho (|i\rangle_A \otimes |2\rangle_B) = \delta_{\mu i} |i\rangle_A$

Defⁿ Partial trace $\text{Tr}_B : \mathcal{L}(\mathcal{H}_{AB}) \rightarrow \mathcal{L}(\mathcal{H}_A)$
 $\text{Tr}_B \rho_{AB} = \sum_{\mu=1}^{d_B} \langle \mu | \rho_{AB} | \mu \rangle_B$

ex $\rho_{AB} = \sum_{i\mu} \sum_{j\nu} \rho_{i\mu j\nu} |i\rangle_A | \mu \rangle_B \langle j | \langle \nu |$

$$\begin{aligned} \text{Tr}_B \rho_{AB} &= \sum_{\lambda} \langle \lambda | \rho_{AB} \left(\sum_{i\mu} \sum_{j\nu} \rho_{i\mu j\nu} |i\rangle_A | \mu \rangle_B \langle j | \langle \nu | \right) | \lambda \rangle_B \\ &= \sum_{i\mu} \sum_{j\nu} \rho_{i\mu j\nu} |i\rangle_A \langle j | \left(\sum_{\lambda} \underbrace{\langle \lambda | \mu \rangle_B}_{\delta_{\lambda\mu}} \underbrace{\langle \nu | \lambda \rangle_B}_{\delta_{\lambda\nu}} \right) \\ &= \sum_{ij} \left(\sum_{\mu} \rho_{i\mu j\mu} \right) |i\rangle_A \langle j | \\ &\equiv \rho_A \quad \text{reduced state on } A \end{aligned}$$

Entropy

Defⁿ Let $\rho \in \mathcal{S}(\mathcal{H})$. The Von Neumann entropy of ρ is

$$S(\rho) = - \text{Tr} \rho \log \rho$$

ex $\rho = \sum_i p_i |i\rangle \langle i|$ then $S(\rho) = - \sum_i p_i \log p_i$

- a measure of purity

- a measure of bipartite entanglement

// in fact, the unique measure of pure-state bipartite entanglement

Properties: (assume $\dim \mathcal{H} = d < \infty$ for now)

1. $\max_{\rho} S(\rho) = \log d$

Proof: work in eigenbasis of ρ : $S(\rho) = - \sum_{i=1}^d p_i \log p_i$

• Maximize $S(\rho) = S(p_1, \dots, p_d)$ subject to
 $0 \leq p_i \leq 1, \quad \sum p_i = 1$

• let $p_d = 1 - \sum_{i=1}^{d-1} p_i$ (take care of last constraint)

• $\frac{\partial S}{\partial p_a} = -\log p_a - 1 - (-\log(1 - \sum_{i=1}^{d-1} p_i) - 1)$

$= \log p_d - \log p_a \stackrel{!}{=} 0$

$\Rightarrow p_d = p_a \quad \forall a = 1, \dots, d-1$

$\Leftrightarrow p_i = \frac{1}{d} \quad \forall i$

$\therefore \max S(\rho) = \log d$, achieved on $\rho = \frac{I}{d}$ maximally mixed

0. $S(\rho) = 0 \Leftrightarrow \rho = |\psi\rangle\langle\psi|$ pure

Proof: \Rightarrow trivial

\Leftarrow above, we found a single critical point of $S(\rho)$, and it was a maximum

\therefore minimum must be achieved at an edge pt. $p_a = 1, p_{i \neq a} = 0$

\Rightarrow this is a pure state, for which $S(\rho) = 0$

$\leadsto S(\rho)$ is a measure of purity

3. Let $|\psi\rangle_{AB} \in \mathcal{H}_{AB}$ be pure, $\rho_A = \text{Tr}_B |\psi\rangle\langle\psi|$ $\rho_B = \text{Tr}_A |\psi\rangle\langle\psi|$.

Then $S(\rho_A) = S(\rho_B)$

Follows from Schmidt decomposition / SVD

4. For $|\psi\rangle_{AB} \in \mathcal{H}_{AB}$, $S(\rho_A) = S(\rho_B) = 0$ iff $|\psi\rangle_{AB} = |\phi\rangle_A \otimes |\chi\rangle_B$,

i.e. $|\psi\rangle_{AB}$ is unentangled across A, B

\rightarrow for bipartite ρ_{AB} , $S(\rho_A), S(\rho_B)$ is sometimes called entanglement entropy

5. $S(U\rho U^\dagger) = S(\rho)$ for any unitary U

$\rightarrow S$ is invariant under local operations

Entropy also obeys a lot of inequalities, e.g.

• Subadditivity: $S(\rho_{AB}) \leq S(\rho_A) + S(\rho_B)$
equality iff $\rho_A = \rho_A \otimes \rho_B$

• Araki-Lieb: $|S(\rho_A) - S(\rho_B)| \leq S(\rho_{AB})$

• Strong Subadditivity: $S(\rho_{AB}) + S(\rho_{BC}) \geq S(\rho_{ABC}) + S(\rho_C)$

... and more!

- there are many more entropic quantities that one can define/interpret
- an important one for us is **relative entropy**

Defⁿ Let $\rho, \sigma \in \mathcal{S}(\mathcal{H})$. The **relative entropy** of ρ and σ is

$$D(\rho \parallel \sigma) = \text{Tr}(\rho \log \rho) - \text{Tr}(\rho \log \sigma)$$

* Key property ①: $D(\rho \parallel \sigma) \geq 0$

$$D(\rho \parallel \sigma) = 0 \quad \text{iff} \quad \rho = \sigma$$

(will prove in exercise session)

* Key property ②: $D(\rho \parallel \sigma) \geq \frac{1}{2 \ln 2} \|\rho - \sigma\|_1^2$

(Pinsker's ineq. - won't prove; see Wilde)

Why is this important?

→ Data-processing inequality (next lecture)

→ $\|\rho - \sigma\|_1$ is a measure of the distinguishability of ρ, σ
(exercise session)

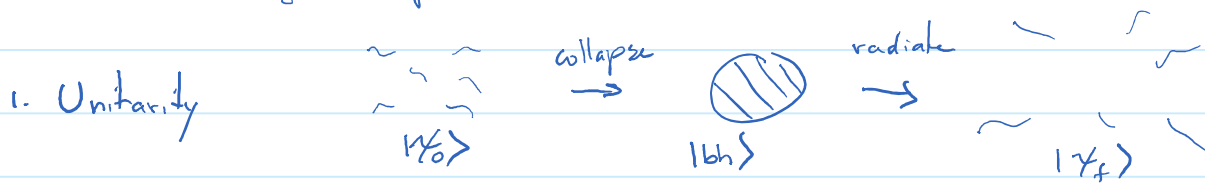
∴ $D(\rho \parallel \sigma)$ small $\Rightarrow \rho, \sigma$ are "close"

Defⁿ $\|\rho\|_1 = \text{Tr}(\sqrt{\rho \rho^\dagger})$

The Black Hole Information Problem

Modern version, due to Almheiri, Marolf, Polchinski, Stanford, Sully:

The following 4 postulates are inconsistent:

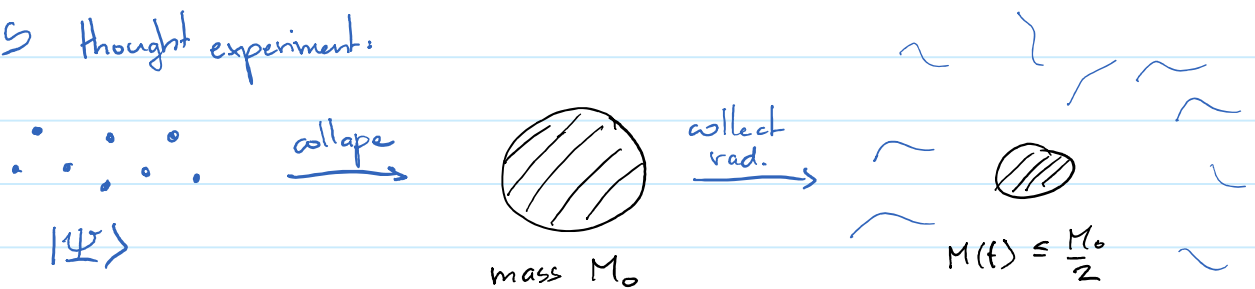


2. L.E.F.T.: outside bh stretched horizon, physics described by local effective field theory

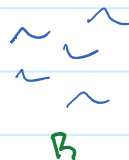
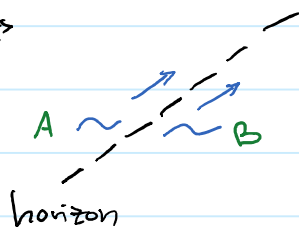
3. Black holes are QM systems: discrete spectrum of states, density given by $S_{BH} = \frac{A}{4G_N}$

4. No Drama: nothing special (beyond semi-class GR + EFT) happens at stretched horizon

AMPS thought experiment:



act
r



A: outgoing mode inside hor.
B: " " " outside "
B: "early" radiation

(i) SSA: $S_{AB} + S_{BR} \geq S_{ABR} + S_B$

(ii) unitarity $\Rightarrow S_{BR} < S_R$ (B purifies R)

(iii) no drama $\Rightarrow S_{AB} = 0$ (AB in pure near-vac. state)
 $\Rightarrow S_{ABR} = S_R$ (Subadditivity)

∞ $S_R + S_B \stackrel{(iii)}{=} S_{ABR} + S_B \stackrel{(i)}{\leq} S_{AB} + S_{BR} \stackrel{(ii)}{<} S_R \quad \#$

∴ need to give something up

(i) Unitarity → bh destroy info

(ii) L.E.F.T → nonlocality

holography

(iii) BH are QM → remnants

→ beyond QM?

(iv) No Drama → firewalls ($S_{AB} \neq 0$)

→ fuzzballs